

IEEE 802.3 Specifications

- 10Base5: 10Mbps, coaxial (thick) cable, bus topology,
 - maximum segment length 500m, segments may be connected with repeaters, up to 4 repeaters between any two stations, up to 2500m network span, and Manchester encoding
- 10Base2: 10Mbps, coaxial (thin) cable, bus, 200m segments
- 10Base-T: 10Mps, UTP, star, Manchester, 100m links
- 10Base-F: 10Mbps, 850 nm fiber, star, On/Off, 2000m links
- 100-Mbps Ethernet (Fast Ethernet), star topology, 100m links, 200m or 400m network spans, uses full-duplex mode with two physical links between nodes
 - —100BASE-TX uses two STPs or Cat. 5 UTPs
 - -100BASE-FX uses two optical fibers
- 1Gbit or 10Gbit Ethernet
 - —10GBASE-L (long): 1310 nm on single-mode fiber, 10 km

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"Taking Turns" MAC protocols

channel partitioning MAC protocols:

- -share channel efficiently at high load
- —inefficient at low load: delay in channel access, 1/N bandwidth allocated even if only 1 active node!

Random access MAC protocols

- —efficient at low load: single node can fully utilize channel
- —high load: collision overhead

"taking turns" protocols

look for best of both worlds!

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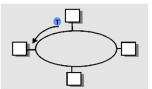
"Taking Turns" MAC protocols

Polling:

- master node "invites" slave nodes to transmit in turn
- Request to Send, Clear to Send msgs
- concerns:
 - polling overhead
 - latency
 - single point of failure (master)

Token passing:

- control token passed from one node to next sequentially.
- token message
- concerns:
 - o token overhead
 - latency
 - single point of failure (token)



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DataLink Layer

5

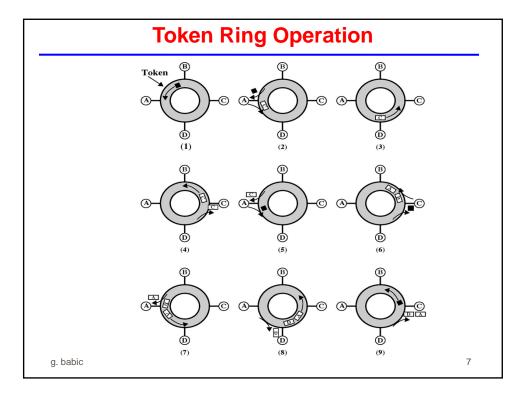
IEEE 802.5 Token Ring

- Small frame (token) circulates when idle
- Station waits for token
- · Changes one bit in token to make it SOF for data frame
- Append rest of data frame
- Frame makes round trip and is absorbed by transmitting station
- Station then inserts new token when transmission has finished.

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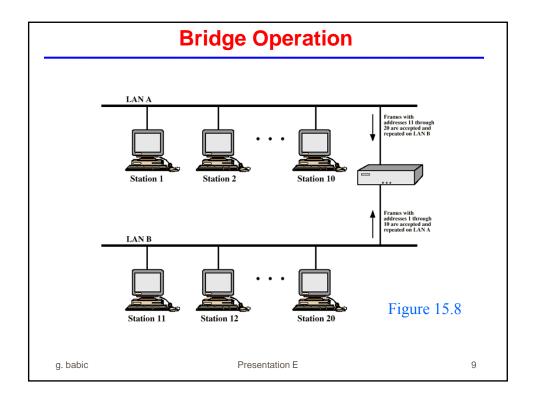
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6



Bridges

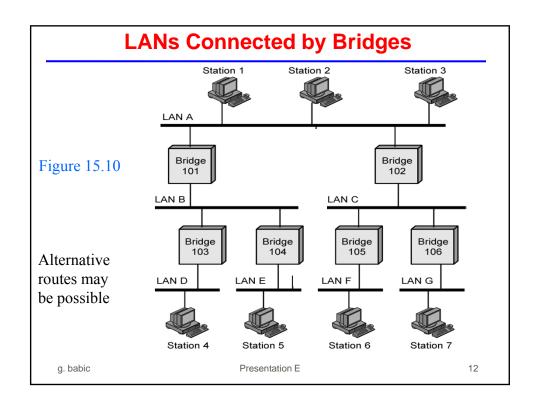
- · Ability to expand beyond single LAN
- Bridge connects two identical (or similar) LANs
 - Identical protocols for physical and link layers
 - Minimal processing
- Router more general purpose
 - Interconnect various LANs and WANs
- · Functioning of bridge
 - Read all frames transmitted on one LAN and accept those address to any station on the other LAN
 - Using MAC protocol for second LAN, retransmit each frame
 - Do the same the other way round



Bridge Design Aspects

- No modification to content or format of frame (for identical LANs)
- No encapsulation
- Exact bitwise copy of frame
- Minimal buffering to meet peak demand
- Bridging is transparent to stations
 - Appears to all stations on multiple LANs as if they are on one single LAN
- Collision does not propagate through bridge
- · Contains routing and address intelligence
 - Must be able to tell which frames to pass
 - May be more than one bridge to cross
- May connect more than two LANs

Bridge Protocol Architecture • IEEE 802.1D • MAC level: station address is at this level • Bridge does not need LLC layer: it is relaying MAC frames User User LLC LLC MAC MAC LAN LAN Physical Physical Physical (a) Architecture t₁, t₈ User Data t_2 , t_7 LLC-H User Data Figure 15.9 t₃, t₄, t₅, t₆ MAC-H LLC-H User Data MAC-T (b) Operation g. babic Presentation E 11



Repeaters, Hubs and Bridges

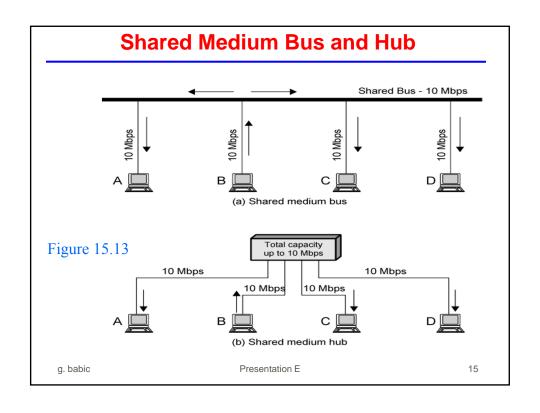
- Repeater: a device that restores data and collision signal
- Hubs: a multi-port repeater with collision detection
- Both a repeater and a hub
 - are physical devices, i.e. they function at the physical layer
 - a main function to extend the length of LAN and/or resolve wiring problem
- Bridges: a device that connects two LANs that use identical or similar protocols at physical layer and MAC sublayer
- Bridge uses MAC sublayer for relaying
 - —<u>Transparent bridge</u>: a bridge that connect LANs with identical MAC protocols
 - —<u>Translation bridge</u>: a bridge that connects LANs with different MAC protocols
 - Encapsulation bridge: two LANs some distance apart can be connected by two bridges that are in turn connected by communication facilities

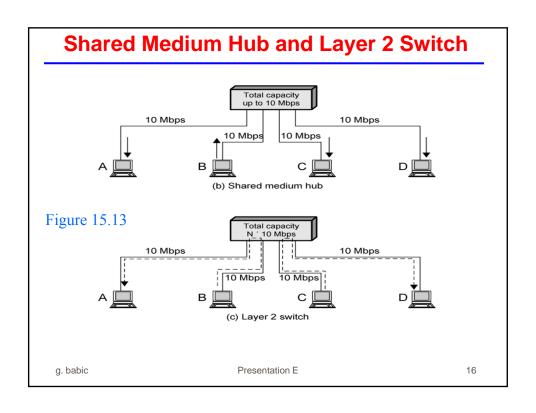
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Hubs

- Active central element of star layout
- Each station connected to hub by two lines (usually two unshielded twisted pairs
 - Transmit and receive
- Limited distances if high data rate and poor transmission qualities of UTP
- Optical fiber may be used for larger distances
- Physically star, logically bus
- Transmission from any station received by all other stations
- Hub acts as a repeater
- When single station transmits, hub repeats signal on outgoing line to each station (that is a header hub; an intermediate hub functions differently)
- If two stations transmit at the same time, collision g. babic Presentation E

14





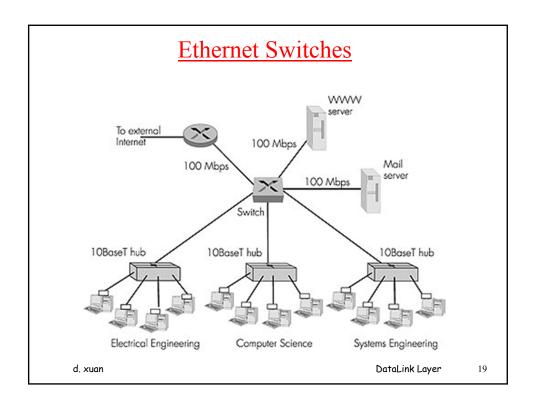
Layer 2 Switches

- Incoming frame from particular station switched to appropriate output line
- Relaying at MAC layer
- · Unused lines can switch other traffic
- More than one station transmitting at a time
- Multiplying capacity of LAN
- No change to attached devices to convert bus LAN or hub LAN to switched LAN, e.g. for Ethernet LAN, each device still uses Ethernet MAC protocol
- Each device has dedicated capacity equal to original LAN, assuming switch has sufficient capacity to keep up with all devices
- A layer 2 switch can be thought as a multi-port bridge with each LAN having only one device; a layer 2 switch selectively forwards frames from one LAN port to another.

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Types of Layer 2 Switches

- Store-and-forward switch
 - Accepts frame on input line
 - Buffers it briefly,
 - Then routes it to appropriate output line
 - Delay between sender and receiver
 - Boosts integrity of network
- Cut-through switch
 - Takes advantage of destination address appearing at beginning of frame
 - Switch begins repeating frame onto output line as soon as it recognizes destination address
 - Highest possible throughput
 - Risk of propagating bad frames, since switch unable to check CRC prior to retransmission



Ring and Star Topologies Usage

- Ring topology
 - Very high speed links over long distances
 - Used in MANs
 - Single link or repeater failure disables network
- Star topology
 - Uses natural layout of wiring in building
 - Best for short distances
 - High data rates for small number of devices

Choice of Topology and Medium

- Choice of topology influenced by:
 - Reliability
 - Expandability
 - Performance
 - Needs considering in context of:
 - Medium
 - Wiring layout
 - Medium access control
- Choice of medium constrained by:
 - LAN topology
 - Capacity
 - Reliability
 - Types of data supported
 - Environmental scope

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Aloha

- Used in packet radio network
- When station has frame, it sends
- Station listens (for maximum round trip time) plus small increment
- If frame OK and address matches receiver, send ACK
- Frame may be damaged by noise or by another station transmitting at the same time (collision)
- Frame check sequence (as in HDLC) detects error
- If ACK received, fine. If no ACK received, retransmit.
- If no ACK after repeated transmissions, give up
- Each station has to have different retransmission time out.
 Why? How?
- Any overlap of frames causes collision
- Max utilization 18% of channel capacity

Slotted Aloha

- Time in uniform slots equal to frame transmission time
- Need central clock (or other sync mechanism)
- Transmission begins at slot boundary
- Everything else as in Aloha
- Frames either miss or overlap totally
- Max utilization 37% of channel capacity

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Values of P?

- If n stations waiting to send, at the end of transmission, expected number of stations attempting to transmit is
 - $-n \times P$
- If n×P > 1 then on average there will be a collision
- Repeated attempts to transmit almost guaranteeing more collisions, since retries compete with new transmissions
- Eventually, all stations trying to send
 - Continuous collisions; zero throughput
- So to void instability n×P < 1 for expected peaks of n
- If heavy load expected, P small
- However, as P made smaller, stations wait longer
- At low loads, this gives very long delays

Binary Exponential Backoff

- In first 10 retransmission attempts, mean value of random delay doubled
- Value then remains same for 6 further attempts
- After 16 unsuccessful attempts, station gives up and reports error
- As congestion increases, stations back off by larger amounts to reduce the probability of collision.
- 1-persistent algorithm with binary exponential backoff efficient over wide range of loads
 - Low loads, 1-persistence guarantees station can seize channel once idle
 - High loads, at least as stable as other techniques
- Backoff algorithm gives last-in, first-out effect
- Stations with few collisions transmit first

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100Mbps Ethernet Specifications

- Traditional Ethernet half duplex
- With full-duplex, station can transmit & receive simultaneously
 - Two physical links between nodes: transmission and reception
- 100-Mbps Ethernet (Fast Ethernet) uses full-duplex mode, theoretical transfer rate 200 Mbps, star topology, 100m links, 200m or 400m network spans
- 100BASE-TX uses two STPs or Cat. 5 UTPs
- 100BASE-FX uses two optical fibers
- 100BASE-T4 uses 4 Cat. 3 (low quality voice-grade) UTP;
 - Taking advantage of large installed base
 - Data stream split into three separate streams; Each with an effective data rate of 33.33 Mbps

10Gbit Ethernet Specifications

- 1Gbit or 10Gbit Ethernet uses full-duplex mode only
- · Fiber links used
- 10GBASE-S (short):
 - 850 nm on multimode fiber and up to 300 m
- 10GBASE-L (long)
 - 1310 nm on single-mode fiber and up to 10 km
- 10GBASE-E (extended)
 - 1550 nm on single-mode fiber and up to 40 km
- 10GBASE-LX4:
 - 1310 nm on single-mode or multimode fiber and up to 10 km
 - Wavelength-division multiplexing (WDM) bit stream across four light waves