

Types in Programming Languages

Chapter 7

Types

- Organization of **untyped values**
 - Untyped universes: bit strings, S-expr, ...
 - Categorize based on usage and behavior
- Type = set of computational entities with uniform behavior
- Constraints to enforce correctness
 - Check the applicability of operations
 - Should not try to multiply two strings
 - Should not use a character value as a condition of an if-statement
 - Should not use an integer as a pointer

Examples of Type Checking

- Built-in operators should get operands of correct types
- Type of left-hand side must agree with the value on the right-hand side
- Procedure calls: check number and type of actual arguments
- Return type should match returned value

Static Typing

- Statically typed languages: expressions in the code have **static types**
 - static type = claim about run-time values
 - Types are either **declared** or **inferred**
 - Examples: C, C++, Java, ML, Pascal, Modula-3
- A statically typed language typically does some form of **static type checking**
 - E.g., at compile time Java checks that the [] operator is applied to a value of type “array”
 - May also do dynamic (run-time) checking
 - e.g., Java checks at run time for array indices out of bounds and for null pointers

Dynamic Typing

- Dynamically-typed languages: entities in the code **do not** have static types
 - Examples: Lisp, Scheme, CLOS, Smalltalk, Perl, Python
 - Entities in the code do not have declared types, and the compiler does not try to infer types for them
- **Dynamic type checking**
 - Before an operation is performed at run time
 - E.g., in Scheme: **(+ 5 #t)** fails at run time, when the evaluation expects to see two numeric atoms as operands of +

Strongly vs. Weakly Typed

- **Strongly typed** languages: type-incorrect operations are not performed at run time
 - Things cannot “go wrong”: no undetected type errors
 - Certain run-time errors are possible but clearly marked as such
 - i.e. array index out of bounds, null pointer
 - C/C++: weakly typed, Java: strongly typed
- Independent of static vs. dynamic
 - Lisp, Scheme, Python: strongly, dynamically typed
 - Forth: weakly, dynamically typed

Examples of Types

- Integers
- Arrays of integers
- Pointers to integers
- Records with fields **int x** and **int y**
 - e.g., “struct” is C
- Objects of class C or a subclass of C
 - e.g., C++, Java, C#
- Functions from any list to integers

Numeric Types

- Varied from language to language
- C does not specify the ranges of numeric types
 - Integer types: char, short, int, long, long long
 - Includes “unsigned” versions of these
 - Floating-point types: float, double, long double
- Java specifies the ranges of numeric types
 - byte: 8-bit signed two's complement integer [-128,+127]
 - short: 16-bit signed two's complement integer [-32768,+32,767]
 - int: 32-bit signed two's complement integer [-2147483648,+2147483647]
 - long: 64-bit signed two's complement integer [-9223372036854775808, +9223372036854775807]
 - float/double: single/double-precision 32-bit IEEE 754 floating point
 - char: single 16-bit Unicode character; minimum value of '\u0000' (or 0) and a maximum value of '\uffff' (or 65535)

Enumeration Types

- C: a set of named integer constant values
 - Example from the C specification

```
enum hue { chartreuse, burgundy, claret=20, winedark };  
/* the set of integer constant values is { 0, 1, 20, 21 } */  
enum hue col, *cp;  
col = claret; cp = &col;  
if (*cp != burgundy) ...
```
- Java: a fixed set of named items (not integers)

```
enum Day { SUNDAY, MONDAY, TUESDAY, WEDNESDAY,  
THURSDAY, FRIDAY, SATURDAY }
```

 - In reality, it is like a class: e.g., it can contain methods

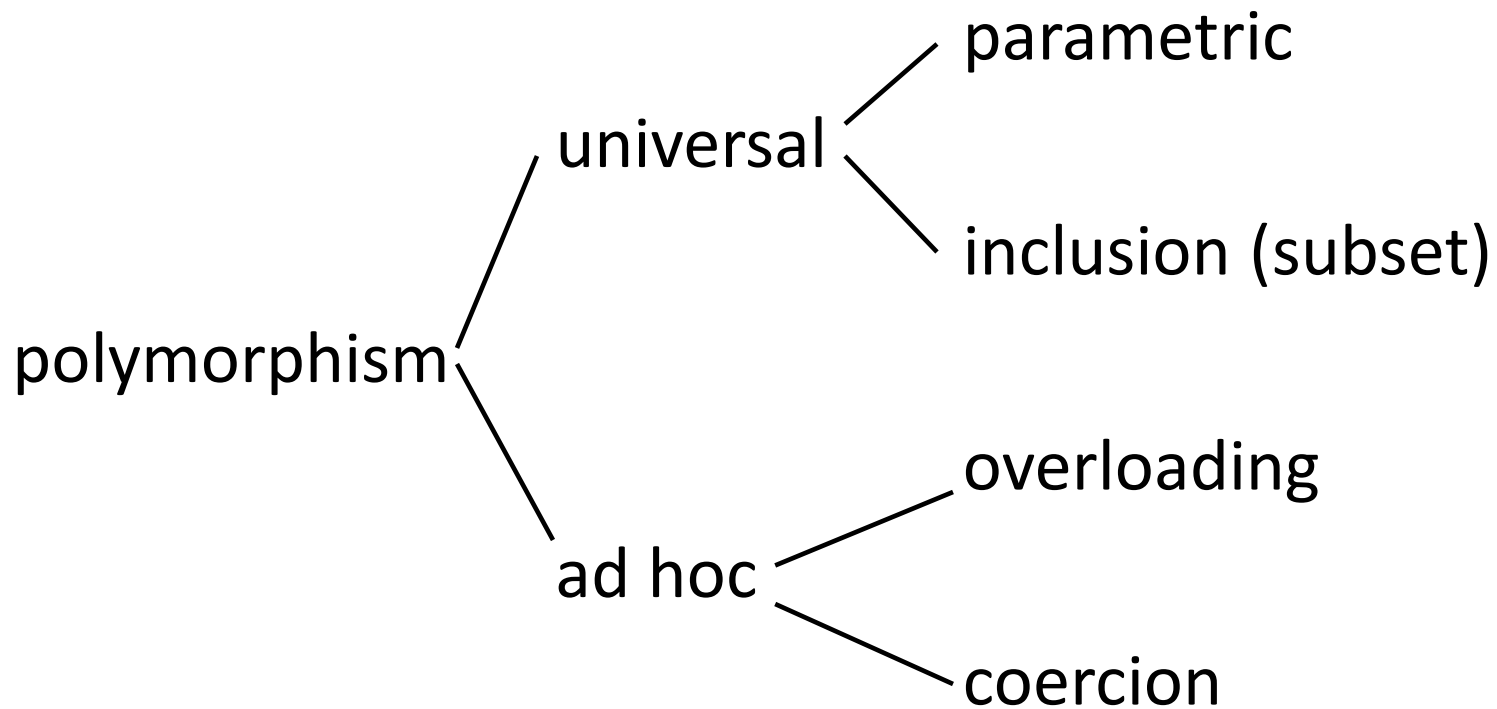
Types as Sets of Values

- Integers
 - Any number that can be represented in 32 bits in signed two's-complement
 - “**type int**” = $\{ -2^{31}, \dots, 2^{31} - 1 \}$
- Class type (not the same as a class)
 - Any object of class C or a subclass of C
 - “**type C**” = set of all instances of C or of any transitive subclass of C (“**class C**” is just a blueprint for objects)
- **Subtypes are subsets**: T2 is a **subtype** of T1 if the T2's set of values is a subset of T1's set of values

Monomorphism vs. Polymorphism

- Greek:
 - mono = single
 - poly = many
 - morph = form
- Monomorphism
 - Every computational entity belongs to exactly one type
- Polymorphism
 - A computational entity can belong to multiple types

Types of Polymorphism



Coercion

- Values of one type are silently converted to another type
 - e.g. addition: $3.0 + 4$: converts 4 to 4.0
 - $\text{int} \times \text{int} \rightarrow \text{int}$ or $\text{real} \times \text{real} \rightarrow \text{real}$
- In a context where the type of an expression is not appropriate
 - either an automatic coercion (conversion) to another type is performed automatically
 - or if not possible: compile-time error

Coercions

- Widening
 - coercing a value into a “larger” type
 - e.g., **int** to **float**, subclass to superclass
- Narrowing
 - coercing a value into a “smaller” type
 - loses information, e.g., **float** to **int**

Widening Primitive Conversions in Java

- Widening primitive conversions
 - byte to short, int, long, float, or double
 - short to int, long, float, or double
 - char to int, long, float, or double
 - int to long, float, or double
 - long to float or double
 - float to double
- “integral type to integral type” and “float to double” do not lose any information

Widening Primitive Conversions in Java

- Language Spec says
 - Conversion of an int or long value to float, or of a long value to double, may result in loss of precision
 - The result may lose some of the least significant bits of the value. In this case, the resulting floating-point value will be a correctly rounded version of the integer value, using IEEE 754 round-to-nearest mode

Contexts for Widening Conversions

- **Assignment conversion**: when the value of an expression is assigned to a variable
- **Method invocation conversion**: applied to each argument value in a method or constructor invocation
 - The type of the argument expression must be converted to the type of the corresponding formal parameter
- **Casting conversion**: applied to the operand of a cast operator: (float) 5

Contexts for Widening Conversions

- **Numeric promotion**: converts operands of a numeric operator to a common type
- Example: binary numeric promotion
 - e.g. +, -, *, etc.
 - If either operand is double, the other is converted to double
 - Otherwise, if either operand is of type float, the other is converted to float
 - Otherwise, if either operand is of type long, the other is converted to long
 - Otherwise, both are converted to type int

Narrowing Conversions

- Narrowing primitive conversions in Java
 - e.g. long to byte, short, char, or int
 - float to byte, short, char, int, or long
 - double to byte, short, char, int, long, or float
- Examples of loss of information
 - int to short loses high bits
 - int not fitting in byte changes sign and magnitude
 - double too small for float underflows to zero

Polymorphism by Overloading

- Multiple definitions of the same name
- E.g. name **+** for several operations: has several types (**function types** $X \rightarrow Y$)
 - $\text{double} \times \text{double} \rightarrow \text{double}$ (binary plus)
 - $\text{float} \times \text{float} \rightarrow \text{float}$
 - $\text{long} \times \text{long} \rightarrow \text{long}$
 - $\text{int} \times \text{int} \rightarrow \text{int}$
 - $\text{double} \rightarrow \text{double}$ (unary plus)
 - $\text{float} \rightarrow \text{float}$
 - $\text{long} \rightarrow \text{long}$
 - $\text{int} \rightarrow \text{int}$

Overloading vs. Overriding in Java

```
class Point {  
    int x = 0, y = 0;  
    void move(int dx, int dy) { x += dx; y += dy; } }  
  
class RealPoint extends Point {  
    float x = 0.0, y = 0.0;  
    void move(int dx, int dy)  
        { move((float)dx, (float)dy); }  
    void move(float dx, float dy)  
        { x += dx; y += dy; } }
```

Overloading vs. Overriding in Java

```
public static void main(String[] args) {  
    RealPoint rp = new RealPoint();  
    // compile-time resolution: the most specific  
    // target method  
    rp.move(1.5f, 1.5f); → RealPoint.move(float,float)  
    rp.move(2,2); → RealPoint.move(int,int)  
    Point p = rp;  
    p.move(3,3); → compile time: Point.move(int,int)  
                 → run time: RealPoint.move(int,int)  
    // can we say p.move(3.3f,3.3f)?  
}
```

Overloading: Most Specific Method

```
class Test {  
    static void test(RealPoint p, Point q) { ... }  
    static void test(Point p, RealPoint q) { ... }  
    public static void main(String[] args) {  
        RealPoint rp = new RealPoint();  
        test(rp,rp); // compile-time error  
    }  
}
```

Parametric Polymorphism: Generics in Ada

generic

type T is private;

function Id(X : in T) return T is

begin

return X;

end;

function IntId is new Id (INTEGER);

function FloatId is new Id (FLOAT);

*The type is the
function type $T \rightarrow T$*

*Similar: templates in C++; generics in Java 1.5 and later
generic functions and classes*

Parametric Polymorphism: Generics in Java

```
package java.util;  
  
public interface Set<E> extends Collection<E> { ...  
    Iterator<E> iterator();  
    boolean add(E e);  
    boolean addAll(Collection<? extends E> c); }  
  
class Rectangle { ... }  
class SwissRectangle extends Rectangle { ... }  
  
Set<Rectangle> s = new HashSet<Rectangle>();  
s.add(new Rectangle(1.,2.)); s.add(new SwissRectangle(3.,4.,5));  
Set<SwissRectangle> s2 = new TreeSet<SwissRectangle>();  
s2.add(new SwissRectangle(6.,7.,8)); s.addAll(s2);
```

Inclusion (Subset) Polymorphism

- Subtype relationships among types
 - Defined by “Y is subset of X” (i.e., set inclusion)
- A computational entity of a subtype may be used in any context that expects an entity of a supertype
- Typical examples
 - Imperative languages: record types
 - Object-oriented languages: class types

Subtyping in Java

- Recall that **class type C** is the set of all instances of class C or of any transitive subclass of C

- Subtyping between class types

```
class X { int m () { ... } }
```

```
class Y extends X { int m () { ... } }
```

```
X x = new Y();
```

```
int i = x.m();
```

- Interface type: the set of all instances of classes that implements the interface (transitively)

```
interface Z { bool m(); }
```

```
class W implements Z { bool m() { ... } }
```

```
Z z = new W(); bool b = z.m();
```

Example: The Core Interpreter

- To illustrate these issues, consider again the implementation of the interpreter
- Use of PT array or PT data type
 - The compiler will not stop us from creating a `<decl>` with a child `<stmt>`, instead of `<id-list>`
 - Conceptually, this is a type error: but our program does not declare “rich enough” types to catch such an error
- Solution: create a separate type for each non-terminal (e.g., a separate class type)

Class **Prog** for Non-Terminal <prog>

```
class Prog {  
    private: DeclSeq* decl_seq; StmtSeq* stmt_seq;  
    public:  
        Prog() { decl_seq = NULL; stmt_seq = NULL; }  
        void parse() {  
            scanner->match(PROGRAM);  
            decl_seq = new DeclSeq(); decl_seq->parse();  
            scanner->match(BEGIN);  
            stmt_seq = new StmtSeq(); stmt_seq->parse();  
            scanner->match(END); scanner->match(EOF);  
        }  
        void print() {  
            cout << "program "; decl_seq->print();  
            cout << " begin "; stmt_seq->print(); cout << " end";  
        }  
        void exec() {  
            decl_seq->exec(); stmt_seq->exec();  
        }  
};
```

Class `StmtSeq` for Non-Terminal `<stmt-seq>`

```
class StmtSeq {
private: Stmt* stmt; StmtSeq* stmt_seq;
public:
  StmtSeq() { stmt = NULL; stmt_seq = NULL; }
  void parse() {
    stmt = new Stmt(); stmt->parse();
    if (scanner->currentToken() == END) return;
    // Same for ELSE, ENDIF, ENDWHILE
    stmt_seq = new StmtSeq(); stmt_seq->parse();
  }
  void print() {
    stmt->print();
    if (stmt_seq != NULL) stmt_seq->print();
  }
  void exec() {
    stmt->exec();
    if (stmt_seq != NULL) stmt_seq->exec();
  }
};
```

Another Example

- $\langle \text{exp} \rangle ::= \text{int} \mid \langle \text{exp} \rangle + \langle \text{exp} \rangle \mid \langle \text{exp} \rangle * \langle \text{exp} \rangle$
 - Not Core, but similar to parts of Core
- For each expression, want to be able to
 - Compute its value: **int evalExpr()**
 - Determine if its value is even: **bool isEven()**
- What classes should we have? What are the methods in these classes? What are the bodies of those methods?