

Dimension-Independent Simplification and Multi-Resolution Representation of Morse Complexes

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Abstract—We have defined and implemented atomic and dimension-independent simplification operators on a graph-based representation of Morse complexes, and we have defined and implemented a multi-resolution model for Morse complexes built through such simplification operators.

1 INTRODUCTION

Morse theory [9] offers a natural and intuitive way of analyzing the structure of a scalar field f and of compactly representing it through decompositions of the domain M of f into meaningful regions associated with its critical points, called *Morse and Morse-Smale complexes*. A descending (ascending, resp.) Morse complex Γ_d (Γ_a) partitions manifold M into cells defined by integral lines converging to (originating at, resp.) critical points of f . For a survey of algorithms for computing Morse and Morse-Smale complexes, see [1, 4]. In our work, we use the algorithm in [10] which is based on discrete Morse theory and has been recently used in the community [5]. Our implementation starts from a decomposition of the domain of the scalar field f into a simplicial complex.

Morse and Morse-Smale complexes built on current large-size data sets can be quite large, and the noise in the data often produces over-segmentations. Simplification of these complexes can be achieved by applying the *cancellation* operator [8]. Cancellation has been used to build hierarchical models for representing Morse complexes in 2D and 3D [7].

2 SIMPLIFICATION OPERATORS ON MORSE COMPLEXES

In [2], we have defined atomic and dimension-independent simplification and refinement operators on Morse complexes, which form a minimally complete set of operators for creating and updating Morse and Morse-Smale complexes. The *remove* simplification operator collapses two saddles of consecutive index that are connected through a unique integral line, and such that one of them is connected to at most one saddle of the same index as the other one. Operator $remove_{i,i+1}(q, p, p')$ applies when i -saddle q is connected to exactly one other $(i+1)$ -saddle p' different from p . It collapses i -saddle q and $(i+1)$ -saddle p into $(i+1)$ -saddle p' . In the descending complex Γ_d , it collapses i -cell q and $(i+1)$ -cell p into a unique $(i+1)$ -cell p' incident in q and different from p . An example of $remove_{1,2}(q, p, p')$ on a 2D descending Morse complex is illustrated in Figure 1: 1-cell q is deleted and 2-cell p is merged into 2-cell p' . Operator $remove_{i,i-1}(q, p, p')$ is completely dual.

We encode the topology of the descending and ascending Morse complexes Γ_a and Γ_d , respectively, through the *Morse Incidence Graph (MIG)* [3] (see Figure 1 for a 2D example). An MIG is a dimension-independent multigraph $G = (N, A)$ such that (i) there is a one-to-one correspondence between the nodes in N and the i -cells of

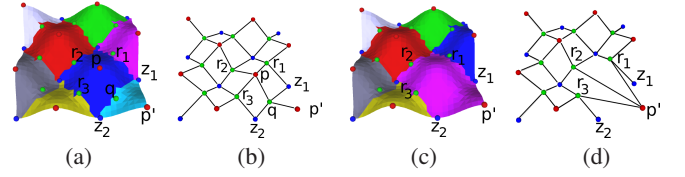


Fig. 1. The descending 2D Morse complex Γ_d for function $f(x, y) = \cos.x\cos.y$ (a) and the corresponding MIG (b). An example of $remove_{1,2}(q, p, p')$ on Γ_d (c), and on the MIG (d). After $remove_{1,2}(q, p, p')$, nodes p and q and the incident arcs are deleted. Arcs (p, r_j) are replaced with arcs (p', r_j) , $1 \leq j \leq 3$.

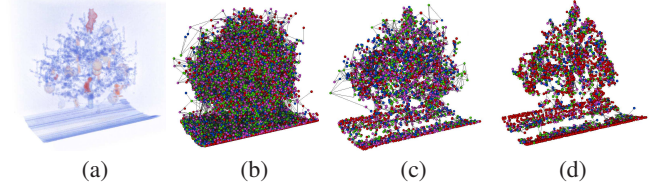


Fig. 2. The field behavior for the *xMaxTree* data set (a). The MIG at full resolution (b), the MIG after 10K (c) and after 16K simplifications (d).

Γ_d , (and thus the $(n-i)$ -cells of Γ_d) (called i -nodes), and (ii) there are k arcs joining an i -node p with an $(i+1)$ -node q if and only if i -cell p appears k times on the boundary of $(i+1)$ -cell q in Γ_d .

A *remove* operator on an MIG $G = (N, A)$ eliminates the two nodes corresponding to the two critical points p and q . A $remove_{i,i+1}(q, p, p')$ is feasible on G if there is a unique arc (p, q) in A , and there is a unique $(i+1)$ -node $p' \in N$ different from $(i+1)$ -node p and connected to i -node q . The effect of the *remove* operator on the MIG is completely local, since it deletes the arcs incident in q and the arcs incident in p and in $(i+2)$ -nodes in the MIG, and it modifies the arcs incident in p and in i -nodes in the MIG which become incident in p' .

The *cancellation* operator [7, 8] eliminates an i -saddle q and an $(i+1)$ -saddle p that are connected through a unique integral line. There is no limit on the number of saddles connected to q and of the same index as the other saddle (i.e., p' is not unique). Let T be the set of $(i+1)$ -nodes connected to q . The effect of the *cancellation* of p and q on the MIG consists of deleting nodes p and q and all arcs incident in either of them, and creating arcs connecting each i -node previously connected by an arc to $(i+1)$ -node p with each node in T .

Thus, a *cancellation* deletes two nodes from the MIG, but, unlike *remove*, it may increase the number of arcs. Some strategies have been proposed to bound the number of arcs introduced by a *cancellation* [6].

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Name	N Simpl	Nodes	Arcs	Cost	Time
<i>Aneurism</i>	-	24082	156098	1.1	-
<i>cancellation</i>	2000	20082	125524	1.03	762.0
<i>remove</i>	2000	20082	97684	0.82	204.17
<i>Bucky</i>	-	2645	9412	0.06	-
<i>cancellation</i>	350	1945	6494	0.05	0.85
<i>remove</i>	350	1945	3586	0.03	0.47
<i>F16</i>	-	32866	146549	1.04	-
<i>cancellation</i>	4000	24866	102847	0.9	300.08
<i>remove</i>	4000	24866	71646	0.6	196.82
<i>Sf2c</i>	-	8826	22778	0.14	-
<i>cancellation</i>	1500	5826	12645	0.12	3.69
<i>remove</i>	1500	5826	9365	0.09	3.18
<i>Sheep</i>	-	83089	684487	4.8	-
<i>cancellation</i>	10000	63089	502714	3.8	18477.3
<i>remove</i>	10000	63089	379491	3.1	5742.32
<i>VisMale</i>	-	16766	64695	0.38	-
<i>cancellation</i>	2000	12766	40501	0.36	36.86
<i>remove</i>	2000	12766	31549	0.29	31.31

Table 1. Comparison between *cancellation* and *remove*.

3 EXPERIMENTAL RESULTS

We have developed a simplification algorithm for Morse complexes including operators on the dimension-independent *MIG* based on the *remove* operators. We associate a value $P(p, q)$ to any feasible remove operator, deleting two critical points p and q , as the absolute difference in function values at p and q , and we perform simplifications in order of increasing $P(p, q)$ values. We have performed experiments on the simplification of Morse complexes by using six data sets describing 2D scalar fields, and eight data sets describing 3D scalar fields on a 3.2GHz processor with 2 GBytes of memory. We have used different threshold values on $P(p, q)$: 1% of the maximum $P(p, q)$ value for light noise removal, 10% for stronger noise removal, and 20% or greater for consistently reducing the complexity of the *MIG*. The storage cost of the simplified *MIG* using these three different thresholds is equal to 95%, 65% and 35% of the cost of the *MIG* at full resolution.

In Table 1, we show the results obtained by comparing the *remove* operator with the *cancellation* operator. For each data set, we show in the first row the number of nodes and arcs in the full resolution *MIG*. In the second and third rows, we show the statistics related to *cancellation* and *remove* operators, respectively: the number of simplifications applied, the number of nodes and arcs in the simplified *MIG*, the cost of the data structure encoding the *MIG* (in Mb), and the time (in sec) needed to perform the simplifications. The number of arcs in the graph simplified with *cancellation* always exceeds the number of arcs in the graph simplified with the same number of *remove*. Such behavior influences the efficiency of the whole algorithm, doubling the time needed to manage and enqueue a larger number of arcs (and thus, a greater number of possible simplifications) for large data sets. The cost of the *MIG* is reduced by 10% ~ 20% by using *remove* instead of *cancellation*.

We have applied our simplification algorithm on a 3D *xMasTree* data set (see Figure 2 (a)), which represents a tomography of a Christmas tree. The *MIGs* at full-resolution, and after 10K and 16K simplifications are shown in Figures 2 (b), (c) and (d), respectively.

4 THE MULTI-RESOLUTION MORSE INCIDENCE GRAPH

We have defined the inverse *insert* refinement operator which undoes the effect of the *remove* simplification operator, and we have designed and implemented a multi-resolution model for the *MIG* representation of the Morse complexes, that we call the *Multi-Resolution Morse Incidence Graph (MMIG)* [3]. An *MMIG* is built from the full-resolution *MIG* by applying a sequence of *remove* operators. It consists of the base *MIG* G_B , representing the Morse complexes at the coarsest resolution, the set of *insert* refinement modifications, inverse to the *remove* applied in the simplification phase, and a dependency relation between refinement modifications, which has been shown to be a partial order relation. From an *MMIG*, it is possible to extract representations of the Morse complexes at any intermediate (uniform or variable) resolu-

tion by considering a subset of refinements that is closed with respect to the partial order.

We have implemented two data structures for storing the *MMIG*, that we call the *implicit* and the *explicit* data structures, respectively. In the *implicit* data structure, the direct dependencies of a modification μ are encoded in two arrays, *ancestors* and *descendants*, and a vector $pq\text{-ancestors}[j]$ encodes whether a modification μ depends on node p , on node q or on both in *ancestors*[j]. Thus, the sets of nodes that will be connected to q or p are implicitly encoded. In the *explicit* data structure, three lists are created for each modification μ , storing the actual pointers to the above mentioned sets of nodes. The *implicit* data structure occupies 3 to 10 times less space than the *explicit* one. The extraction algorithm is from 33% to 50% faster when using the *explicit* data structure.

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