Bounding Volumes


## Bounding Volumes

Use simple volume
enclose object(s)
if ray doesn't intersect volume
it doesn't intersect what's inside
tradeoff for rays where there is
extra intersection test for
object intersections
volume intersections, but not object intersections v.
quick test for no intersection for no volume intersection

## Bounding Volumes

## 3 approaches:

Bound object

Bound screen area that object projects to

Bound area of world space
Can use hierarchical organization of bounding volumes

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## Bound Object

Easy-to-compute approximation to object

Easy to test for ray-bounding-object intersection

Trade-off complexity of computation v. tightness of fit
Can bound object in object space or world space


## Axis-Aligned Bounding box

$\operatorname{limit}[3][0]=10000000000 ; \operatorname{limit}[3][1]=-1000000000$;
for each point, for each dimensions
if $\mathrm{p}[\mathrm{i}][\mathrm{j}]<\operatorname{limit}[j][0]$ then limitIndex[j][0] = i;
if $\mathrm{p}[\mathrm{i}][\mathrm{j}]>\operatorname{limit}[\mathrm{j}][1]$ then limitIndex[j][1] = i;


## Bounding Spheres

Takes some effort to compute optimal bounding sphere
Easy to compute approximation (non-optimal fit)
Easy to test for intersection (ray-sphere)
Not tight fitting for some objects

## Bounding Spheres

Loop through points and record $\mathrm{min} / \mathrm{max}$ in $\mathrm{x}, \mathrm{y}, \mathrm{z}$
Use maximally separated pair of points and their midpoint as initial approximation to sphere

For each point in original set, adjust the
bounding sphere to include the point


## Bounding Spheres

limit[3][0] = -10000000000; limit[3][1] = 1000000000; for each point, for each dimensions
if $\mathrm{p}[\mathrm{i}][\mathrm{j}]<\operatorname{limit}[j][0]$ then limitIndex[j][0] = i;
if p[i][j] > limit[j][1] then limitIndex[j][1] = i;


## Bounding Spheres

$\mathrm{k}=0$;
if $(|\operatorname{limit}[1][1]-\operatorname{limit}[1][0]|>|\operatorname{limit}[0][1]-\operatorname{limit}[0][0]|) \mathrm{k}=1$; if ( $\mid$ limit[2][1]-limit[2][0]| > |limit[k][1]-limit[k][0]| ) k=2; midpoint $=(\mathrm{p}[$ limit $[\mathrm{k}]]+\mathrm{p}[$ limit[k]])/2;
radius $=(\mathrm{p}[$ limit[k]]-p[limit[k]]])/2;


## Bounding <br> Slabs

- For each point
-if p[i] is outside of radius of midpoint
- radius $=($ radius + dist(midpoint, $p[i])) / 2$
- center = p[i]+radius*(center-p[i])/|center-p[i]|;


For each slab, user defines normal to use for slab pair

For each object, compute 2 d's for each N


Takes some effort to compute d's - how?

Takes some effort to test for intersection - how?
Can add more slabs to get tighter fit


## Bounding Slabs

Slab defined by: N, d1, d2

## Convex Hull

Smallest convex polyhedron containing object (point set)


- compute
- t1 = (d1-N.P)/(N.D); t2 = (d1-N.P)/(N.D)
- keep track of entering max, exiting min
- how to determine entering, exiting status?

[^0]- ray defined by $\mathrm{P}(\mathrm{t})=\mathrm{P}+\mathrm{tD}$
- intersection with plane: N.P + t(N.D) = d
- for each slab
- retrieve d1, d2, N


## Convex Hull

Find highest vertex
Find plane through vertex parallel to ground plane
Find second vertex that makes minimum angle with first vertex and up vector


## Convex Hull

Find third vertex that makes plane whose normal makes minimum angle with up vector


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## Convex Hull

In the final convex hull, each edge will be shared by two
and only two triangles
For each unmatched edge (until there are no more),
find vertex that, when a triangle is formed with
the edge, will minimize angle between its normal and
normal of shared face


Hierarchical Bounding Volumes
Compute bounding volume for groups of objects


Compute bounding volume for groups of groups of objects

Test higher-level
bound volumes first

## Bound Area of Projection

Project object to picture plane

Bound colums and rows that object projects to
Only intersect first-level rays with pixels in projected area i.e., only good for ray-casting part of ray-tracing

## Bound Area of Projection



Project vertices onto picture plane

Find 2D bouding box on picture plane


## Bound Area of Projection

Areas of projection can be grouped hierarchically


## Bound Area, not Object

In 2D - it looks like this:

|  |  |  |  |
| :--- | :--- | :--- | :--- |
|  |  |  |  |
|  |  | 0 |  |
|  |  |  | 8 |
|  |  |  |  |

Bucket sort objects into cells

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## Bound Area of 3D World



Divide world space into cells
Dump objects into cells: an object is dumped into each cell it touches

Bound Area of World


Trace ray through cells from closest to farthest Intersect ray with each object in cell Stop when it hits closest object in cell

## Bound Area of World



Traverse cells in
order - test all
objects in cell
If no intersections, step to next cell

If one or more intersections, get
closest intersection

Bound Area of World


Hierarchical approach: cells, then subcells..
Takes significant coding to keep track of level
Overhead in popping up and down in hierarchy
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Binary Spatial Partitioning



[^0]:    Takes some effort to test for intersection - how?

